

Parallel Algorithms and Parallel Computers

(in4026)

Examination June 28, 2010

14.00-17.00 hrs

This exam consists of two parts: Part I (architecture) and Part II (algorithms). If you have done Part I (April 2010), you can do only Part II or both Part I and Part II. If you have not done Part I in April 2010, you have to do both parts.

Please do note that after the August 2010 re-examination, the individual results for Part I and Part II obtained in 2010 or earlier are no longer valid!

IMPORTANT: Please use a separate sheet for each question!

PART I

Question 1 (25 points)

- a. Given is a multistage interconnection network with p processors. Every switching element in the network consists of a 8 by 8 crossbar switch. The interconnection network has 2 stages. How many processors can be connected to the multistage network to make an arbitrary connection between the processors possible?
- b. Give three arguments for selecting a network with a mesh topology.
- c. Given is a 2D mesh without wrap-around connections met p processors. Give expressions for the diameter, the bisection width, the arc connectivity, and the number of links.
- d. Given is a ring with p processors. What is the maximal communication time of a message of length m if we use *Cut-Through* routing.
- e. Define the ideal model of a parallel computer. Give four subclasses of this model with their properties. Which subclass is the most powerful one and why?

Question 2 (25 points)

- a. Show how to embed a p -node 3D-dimensional Mesh into a p -node Hypercube. What are the allowable values of the mesh dimensions for your embedding
- b. Given is a 2D mesh without wrap around connections. Suppose $P_{Sx, Sy}$ denotes the source processor and $P_{Dx, Dy}$ the destination processor.
 - Explain the working of the XY routing algorithm in a 2D mesh.
 - What is the length of the path as function of Sx , Sy , Dx and Dy ?
 - Is this routing algorithm minimal?
- c. Given is a p processor ring. What is the communication time of a message of length m for a *one-to-all* broadcast using *store-and-forward* routing.
- d. Given is a problem with workload W that has a sequential component W_s . The parallel component W_p is totally parallelizable. Proof that W/W_s is an upperbound to the speedup S .
- e. Given is that the cost of a parallel algorithm and its workload are $C = O(n \log n)$ and $W = \log n$, respectively. Is this parallel algorithm Cost Optimal? Explain your answer.

Part II

Question 3 (20 points)

Let x be an integer and let $A[1..n]$ be an arbitrary array of integers.

a. (5 points) Give a fast parallel algorithm on a CRCW to check whether x occurs as one of the integers in A . What is the work-time complexity of your algorithm?

b. (5 points) Give a $T(n) = O(\log n)$ and $W(n) = O(n)$ algorithm on a CREW that, given x and A , computes $\text{rank}[x:A]$.
Notice that it is **not** required that A is sorted.

c. (10 points) Suppose that it is known that A is a sorted array.
Provide a parallel algorithm on a CREW that computes $\text{rank}[x:A]$ in $o(\log n)$ time and $O(n)$ work.

strictly less than $O(n)$

Question 4 (10 points)

Consider the following (third-order) linear recurrence

$$\begin{aligned}y_0 &= 1 \\y_1 &= 2 \\y_2 &= 2 \\y_i &= y_{i-1} + 3y_{i-2} - 2y_{i-3} \text{ for } i = 3, \dots, n-1.\end{aligned}$$

Let $y = [y_0, \dots, y_{n-1}]$.

a. (5 points) Show how this linear recurrence can be reduced to a first-order linear recurrence.

b. (5 points) Show that there exists a parallel algorithm to compute y in time $T(n) = O(\log n)$. What is the work-complexity of your algorithm?

Question 5 (20 points)

a. (5 points) Give the time-work complexity of the doubly-logarithmic tree algorithm for computing a maximum of an array A of n integers. Why is this algorithm not weakly optimal?

b. (5 points) There is an algorithm that is able to reduce the problem of finding the maximum in an array by reducing it to finding the maximum of an array of $n/(\log \log n)$ numbers without affecting the solution. The time complexity of this algorithm is $O(\log \log n)$ and its work complexity $O(n)$.
Explain how such an algorithm could be obtained by partitioning the array A . You only have to explain the idea of the algorithm.

c. (10 points) Give a combination of this algorithm and the doubly-logarithmic tree algorithm to obtain an algorithm with $O(\log \log n)$ time complexity and $O(n)$ work complexity and justify its complexity.

End of Examination

