

**Deeltentamen / Partial Exam**  
**TW1010 Wiskundige Structuren / Mathematical Structures**  
**Maandag 6 november 2017 / Monday November 6, 2017**  
**9.00-11.00**

Naam / Name: |

Studienummer / Student number: | | | |

Rekenmachines zijn niet toegestaan. Vul de antwoorden in de vakken in. Het cijfer is  $(\text{score}+6)/6$ .  
No calculators allowed. Write the solutions in the fields provided. The grade is  $(\text{score}+6)/6$ .

- 1 Bewijs dat  $13^n - 6^n$  een veelvoud is van 7 voor alle  $n \in \mathbb{N}$ .  
 Prove that  $13^n - 6^n$  is a multiple of 7 for all  $n \in \mathbb{N}$ .

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We'll prove the theorem using induction. For  $n = 1$  we see that  $13^1 - 6^1 = 7$  is indeed a multiple of 7. Now suppose  $13^k - 6^k$  is a multiple of 7. Then we have

$$13^{k+1} - 6^{k+1} = 13(13^k - 6^k) + 7 \cdot 6^k$$

Thus  $13^{k+1} - 6^{k+1}$  is the sum of two multiples of 7, and is itself a multiple of 7 as well.

Using induction we have now shown that  $13^n - 6^n$  is a multiple of 7 for all  $n \in \mathbb{N}$ .

- 2 Laat  $R$  een relatie zijn op een verzameling  $A$ . We nemen aan dat  $R$  symmetrisch en transitief is. Vind de fout in het volgende bewijs. Het is niet de bedoeling dat je laat zien waarom de bewering fout is, maar dat je aangeeft welke zin of welk stukje zin niet klopt, en waarom. 4

Bewering:  $aRa$  voor alle  $a \in A$  (d.w.z.  $R$  is reflexief).

Bewijs: Stel dat  $aRb$ . Dan geldt er ook  $bRa$ , want  $R$  is symmetrisch. Nu geldt er  $aRb$  en  $bRa$ , dus uit transitiviteit van  $R$  volgt dat  $aRa$  voor alle  $a \in A$ .

Let  $R$  be a relation on a set  $A$ . Assume  $R$  is symmetric and transitive.  
Find the error in the following proof. You are not meant to show the statement is false, but to indicate which sentence or part of a sentence is wrong and why.

Statement:  $aRa$  for all  $a \in A$  (i.e.  $R$  is reflexive).

Proof: Suppose  $aRb$ . Then also  $bRa$ , since  $R$  is symmetric. Then  $aRb$  and  $bRa$ , so using transitivity it follows that  $aRa$  for all  $a \in A$ .

“Suppose  $aRb$ .” is the wrong sentence. It is not necessary that, for an arbitrary  $a$ , there exists a  $b$  such that  $aRb$ .

NB: To be completely rigorous, the proof should have started with “Let  $a \in A$  be arbitrary”, but this sentence is often omitted and is not reason the proof is wrong.

- 3a Zij  $f : \mathbb{N} \rightarrow \mathbb{N}$ , en  $C_1, C_2$  zijn deelverzamelingen van  $\mathbb{N}$ . 4  
Geef een tegenvoorbeeld voor de volgende bewering: Als  $f(C_1) = f(C_2)$ , dan  $C_1 = C_2$ .

Let  $f : \mathbb{N} \rightarrow \mathbb{N}$ , and let  $C_1, C_2$  be subsets of  $\mathbb{N}$ .

Give a counterexample to the following statement: If  $f(C_1) = f(C_2)$ , then  $C_1 = C_2$ .

Take  $f(n) = 3$  (the constant function),  $C_1 = \{1\}$  and  $C_2 = \{2\}$ .

Then  $f(C_1) = \{3\} = f(C_2)$  while  $C_1 \neq C_2$ .

3b Maak de definitie af. De functie  $g : A \rightarrow B$  is *injectief* als ...

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Complete the definition. The function  $g : A \rightarrow B$  is *injective* if ...

for all  $a_1, a_2 \in A$  we have  $g(a_1) = g(a_2) \Rightarrow a_1 = a_2$

3c Neem aan dat  $g : A \rightarrow B$  injectief is, en dat  $C_1, C_2$  deelverzamelingen zijn van  $A$ .

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Bewijs: Als  $g(C_1) = g(C_2)$ , dan  $C_1 = C_2$ .

Assume  $g : A \rightarrow B$  is injective, and  $C_1, C_2$  are subsets of  $A$ .

Prove: If  $g(C_1) = g(C_2)$ , then  $C_1 = C_2$ .

Suppose  $g(C_1) = g(C_2)$ . We'll show  $C_1 \subseteq C_2$ . Let  $a \in C_1$ . Then  $g(a) \in g(C_1)$ , so  $g(a) \in g(C_2)$ . Thus there is an  $a_2 \in C_2$  with  $g(a_2) = g(a)$ . By injectivity of  $g$  we find  $a_2 = a$ . Thus we have  $a = a_2 \in C_2$ , and we can conclude  $C_1 \subseteq C_2$ .

By symmetry we likewise obtain  $C_2 \subseteq C_1$ , thus we can conclude  $C_1 = C_2$ .

Under the assumption that  $g$  is injective we have shown the desired implication.

- 4a Formuleer het *welordningsaxioma* voor  $\mathbb{N}$ .  
Formulate the *well-ordering property* of  $\mathbb{N}$ .

2

Every non-empty subset  $S \subseteq \mathbb{N}$  has a smallest element.

- 4b Zij  $x, y \in \mathbb{R}$  met  $x < y$ . Laat zien dat er een rationaal getal  $r$  bestaat, zo dat  $x < r < y$ .  
Let  $x, y \in \mathbb{R}$  with  $x < y$ . Show that there exists a rational number  $r$  such that  $x < r < y$ .

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First suppose  $x > 0$ . By the Archimedean property there is a natural number  $n > 1/(y - x)$ . Let  $q$  be the smallest natural number bigger than  $nx$  (exists due to the Archimedean property and the well-ordering property).

Then  $x < q/n$ , and  $y = x + (y - x) > x + 1/n > q/n$ .

Thus  $r = q/n \in \mathbb{Q}$  satisfies the condition.

In the case  $x \leq 0$ , choose  $k \in \mathbb{N}$  with  $k > -x$  (again Archimedean property).

Then  $y + k > x + k > 0$ , so there exists  $p \in \mathbb{Q}$  with  $x + k < p < y + k$ .

But then  $r = p - k \in \mathbb{Q}$  satisfies  $x < r < y$ .

- 5a Geef de definitie van een *verdichtingspunt* van een verzameling  $S \subseteq \mathbb{R}$ .  
Give the definition of an *accumulation point* of a set  $S \subseteq \mathbb{R}$ .

2

An accumulation point of a set  $S$  is a point  $s \in \mathbb{R}$  such that for every  $\epsilon > 0$  we have  
 $N^*(s; \epsilon) \cap S \neq \emptyset$ .

- 5b  $S \subseteq \mathbb{R}$  is een niet-lege begrensde verzameling met  $x = \sup S$ .  
Bewijs: als  $x \notin S$ , dan is  $x$  een verdichtingspunt van  $S$ .

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Let  $S \subseteq \mathbb{R}$  be a non-empty bounded set with  $x = \sup S$ .  
Prove: if  $x \notin S$ , then  $x$  is an accumulation point of  $S$ .

Suppose  $x = \sup(S) \notin S$ . Let  $\epsilon > 0$ . By definition of supremum for any  $m < x$  there exists  $s \in S$  with  $m < s$ . Thus in particular (with  $m = x - \epsilon$ ) there exists  $s \in S$  with  $s > x - \epsilon$ . By definition of supremum we also have  $s \leq x$ , and as  $s \in S$  and  $x \notin S$  we must have  $s < x$ . Therefore  $s \in N^*(x; \epsilon)$ . Thus we have shown that  $x \in S'$ .

- 6a  $(s_n)$  is een rij en  $s \in \mathbb{R}$ . Geef de definitie van  $\lim s_n = s$ .  
 $(s_n)$  is a sequence and  $s \in \mathbb{R}$ . Give the definition of  $\lim s_n = s$ .

2

$\lim s_n = s$  if for all  $\epsilon > 0$  there exists  $N$  such that for all  $n > N$  we have  $|s - s_n| < \epsilon$ .

- 6b Bepaal de limiet / Determine the limit

2

$$\lim \frac{2n^2 - 4n + 5}{3n^2 + 6}$$

U mag de rekenregels voor limieten en de uitkomst van  $\lim_{n \rightarrow \infty} n^k$  (voor een vaste  $k$ ) gebruiken.

You can use the rules of calculation for limits and the result of  $\lim_{n \rightarrow \infty} n^k$  (for a fixed  $k$ ).

$$\lim \frac{2n^2 - 4n + 5}{3n^2 + 6} = \lim \frac{2 - \frac{4}{n} + \frac{5}{n^2}}{3 + \frac{6}{n^2}} = \frac{2 \lim 1 - 4 \lim \frac{1}{n} + 5 \lim \frac{1}{n^2}}{3 \lim 1 + 6 \lim \frac{1}{n^2}} = \frac{2 - 4 \cdot 0 + 5 \cdot 0}{3 + 6 \cdot 0} = \frac{2}{3}$$

You can apply the rules of calculation as the sequences in the final expression involving limits all converge, and we don't have to divide by zero.

$$\lim \frac{1}{n^2 - 5} = 0.$$

Let  $\epsilon > 0$ . Set  $N = \max(\sqrt{10}, 1/\sqrt{\epsilon/2})$ . Let  $n > N$  be arbitrary. Then  $n^2 - 5 > \frac{1}{2}n^2$  (as  $n > \sqrt{10}$  and therefore  $n^2 > 10$ ). Moreover we have

$$\left| \frac{1}{n^2 - 5} - 0 \right| = \frac{1}{n^2 - 5} < \frac{1}{n^2/2} = \frac{2}{n^2} < \frac{2}{N^2} \leq \epsilon$$

Thus we have shown that  $\lim \frac{1}{n^2 - 5} = 0$ .

Opgave / Exercise  voortgezet (extra ruimte) / continued (extra space)

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