

Exam OPTIMIZATION (AM2020)

19-01-2021

13:30 – 16:00

This exam consists of 6 questions on 3 pages. You can earn 60 points in total. Your grade is determined by dividing the total number of obtained points by 6. Success!

- You only earn points when you use the right method and describe all steps and arguments clearly.
- You may use a calculator, the lecture notes and your own notes. You are **not allowed to collaborate** or to use any other sources.
- Include the following sentence: "This exam is solely undertaken by myself, without any assistance from others." followed by your name.
- Scan your solutions and your student ID and combine the scans into a **single pdf file**. Use a scanner or a scanner app.
- Upload the combined pdf in the assignment folder "Exam" on Brightspace before 16:15 (or 16:40 if you have extra time).
- During the exam, questions can be asked via email to l.j.j.vaniersel@tudelft.nl.
- After the exam, you may be invited for a **face-to-face check** on Wednesday 20-01-2021 between 14:45 and 16:45. Keep this time free and check your student email regularly.

1. (14 points) Indicate for each of the statements below whether they are true (T) or false (F). For each correct answer you get 2 points, for each incorrect answer you get -1 points and for not filling in an answer you get 0 points.

- (i) If an LP has a feasible solution, then it has an optimal solution.
- (ii) A Steepest Descent sequence always converges to a global minimizer.
- (iii) If matrices A and B are both totally unimodular then the matrix $\begin{bmatrix} A & B \end{bmatrix}$ is also totally unimodular.
- (iv) A 3-approximation algorithm for a minimization problem can return a solution with value 16 for an instance with optimal value 8.
- (v) The matrix below is positive definite.

$$A = \begin{bmatrix} 2 & -1 & 0 \\ -1 & 1 & 0 \\ 0 & 0 & 3 \end{bmatrix}$$

(vi) If \mathbf{x}^* is a critical point of a convex function $f(\mathbf{x})$ with continuous first and second partial derivatives, then \mathbf{x}^* is a global minimizer for $f(\mathbf{x})$.

(vii) The function $f(x_1, x_2) = (x_1 + 2x_2)^4$ is coercive.

2. (5 points) **Prove**, without using the Simplex algorithm, that the following LP has a bounded optimum.

$$\begin{array}{ll} \min & z = x_1 - x_2 + 2x_3 \\ \text{s.t.} & 3x_1 - 3x_2 + 4x_3 \leq 5 \\ & -x_1 + 2x_2 - 2x_3 \leq 7 \\ & x_1, x_2, x_3 \geq 0 \end{array}$$

3. (10 points) Consider the problem of scheduling E exams when only a limited number L of students is allowed to be on campus simultaneously. For each exam, we need to decide whether it takes place on-line or on-campus. If it takes place on-campus, you also need to select a time-slot. There are T (disjoint) time-slots. There are s_e students who signed up for exam e . Assume that, if an exam is on-campus, all students who signed up make the exam on-campus. Similarly, if an exam is on-line, all students who signed up make the exam on-line. Finally, assume that each student signed up for exactly one exam. The goal is that a maximum number of students make their exam on-campus. Give an **ILP formulation** of this problem.
4. (10 points) **Prove** that the following problem is NP-complete.

HITTING

Given: a set U , a collection $S = \{S_1, \dots, S_r\}$ of subsets of U (so $S_i \subseteq U$ for $i = 1, \dots, r$) and an integer k

Decide: can we select at most k elements of U such that each of S_1, \dots, S_r contains at least one selected element?

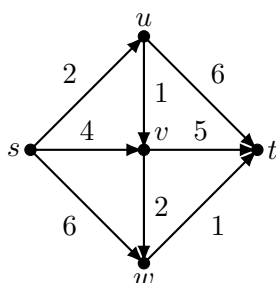
You may use that the following problem is NP-hard. Here, $N^+(v)$ is the set containing vertex v and all its neighbours.

DOMINATION

Given: a graph $G = (V, E)$ and an integer k'

Decide: can we select at most k' vertices from V such that, for each $v \in V$, at least one element of $N^+(v)$ is selected?

5. Consider the following directed graph in which the label of each arc a indicates its length.



- (a) (5 points) Find the **length** of a shortest s - t path using the algorithm of **Dijkstra**.
- (b) (5 points) Use **Complementary Slackness** to **find** a shortest s - t path and to **prove** that no shorter s - t path exists.

- (c) (5 points) Now interpret the labels of the arcs as capacities and **prove**, using a theorem discussed in this course, that no s - t flow of value 8 exists.
6. (6 points) Consider the following Simplex tableau of the LP relaxation of a maximization ILP.

basis	\bar{b}	x_1	x_2	x_3	s_1	s_2
s_1	$17/2$	$-5/2$	0	2	1	$-1/3$
x_2	$7/2$	$1/2$	1	-2	0	$-1/2$
$-z$	$7/2$	$-1/2$	0	0	0	$-1/2$

Find the **Gomory cutting plane** corresponding to the first row of the tableau. You do not need to express it in the original variables x_1, x_2, x_3 .